

CS 61A Discussion 8: **Scheme**

March 23, 2017

Announcements

- Ants is due today. Finish it! Also turn it in!
- HW 6 is due tomorrow. Finish it! Also turn it in!
- HW party today from 6:30-8:30p in 247 Cory.
- Midterm scores are out. Good job everyone!
(Submit regrades by Sunday 4/2 or forever hold your peace.)

116 Attendance: cinnamon_twist
140 Attendance: vanilla_bark



Scheme

As long as 61A lives, it lives
(also I don't know who these people are, so don't ask)

First of all, why are we even learning this
<*insert choice words here*> language?

- To get experience with different languages and paradigms (**functional programming!**)
- So you guys can write an interpreter for it (project 4 hype)



All right, I'm sold. So how do we learn this
<*insert choice words here*> language?

- Test code in Jack's interpreter (scheme.cs61a.org)
- Read Kevin's guide (tiny.cc/kevin_scheme)
- Check out the Scheme reference (tiny.cc/scheme_ref)
- Gaze at infographics (tiny.cc/scheme_illu)
- *Write your own code!* The more you write, the better you will become



Important point for the future

- Everything in Scheme is either a **primitive** or a **combination**.
- Primitives (atomic elements) look like numbers and symbols: `3`, `'4`, ...
- Combinations look like Scheme lists: `(define x 6)`
- Note: in a meta sense, `(define x 6)` is really just a well-formed list containing the elements `'define`, `'x`, and `6`!



A couple of random notes

...that it won't hurt to remember

- The only false values in Scheme are **#f**, **False**, and **false**. Everything else is true!
- Symbols are immutable strings; think of them like variable names or the code itself. They are case-insensitive.
- To get a symbol, use the quote operator:

```
scm> (define x 6)
x
scm> 'x ; the symbol x
x
scm> x ; VALUE of the symbol x
6
```


Definitions

- `define` binds a value to a name (just like `=` in Python). Note that `define` always returns the symbol that has just been assigned a value!
- `(define x 6) ; variable`
- `(define (identity x) x) ; procedure`

Call Expressions

- (`<procedure>` `<arguments>`)
- **Evaluation** (same as Python):
 - Evaluate the operator, then evaluate each of the operands.
 - Apply the operator to the evaluated operands.

Testing for equality

- `=` is used for numbers (and numbers only!)
- `equal?` is used for everything else (...although it actually works for numbers as well)

Special Forms

- `define`, `if`, `and`, `or`,
`not`, `lambda`, `let`
- Special forms look like function calls (**because they're surrounded by parentheses**), but don't follow the same execution process



Lambdas

- `(lambda (<params>) <expression>)`
- Much like Python, lambdas are **first-class function values** and you create new frames when you call them.
- Much like Python, `<expression>` isn't evaluated until the lambda is called.

WWSP?

```
((lambda (x) (x x)) (lambda (y) 4))
```

WWSP?

```
((lambda (x) (x x)) (lambda (y) 4))
```

Answer: 4

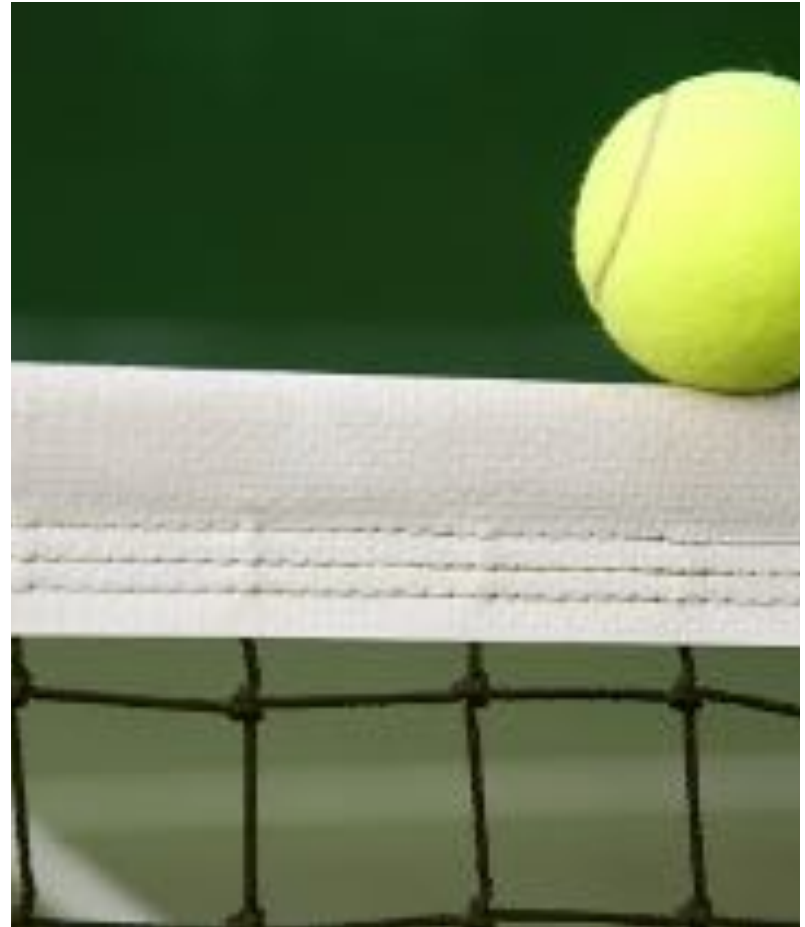
Let

```
(let ((<symbol-1> <value-1>)
      ...
      (<symbol-n> <value-n>))
  <body>)
```

is equivalent to

```
((lambda (<symbol-1> ... <symbol-n>)
  <body>)
  <value-1> ... <value-n>)
```

It's basically saying “assign *these* variables to *these* values, and then execute this code with those assignments in effect.”



Pairs

- `(cons <elt1> <elt2>)` creates a pair containing the elements `<elt1>` and `<elt2>`.
- `(car pair)` selects the first element of a pair.
- `(cdr pair)` selects the second element of a pair.
- `nil`, `()`, `'()` are equivalent and represent the empty list.

Well-formed lists

- A **well-formed list** is a sequence of pairs where the second element of each pair is ALWAYS either another pair or `nil`.
- **Malformed list**: a sequence of pairs where the second element of ANY of those pairs is something other than another pair or `nil`.

Connecting the dots

- The dot delimits the **first** and **second** element of a pair.
- Note that well-formed lists do not have dots in their final interpreter output.
- Rules for displaying a pair in interpreter format:
 - Use a dot to separate the first and second elements of a pair.
 - If the second element is also a pair (i.e. the dot is immediately followed by an open parenthesis), then remove the dot and the next set of parentheses.

- In this way, `(cons 1 (cons 1 2))` becomes `(1 . (1 . 2))` and finally `(1 1 . 2)` when we break it down into the interpreter's final output.

The `list` operator

- `(list <args>)` takes zero or more arguments and returns a **well-formed list** of its arguments (i.e. each argument is in the `car` field of its respective pair).
- `(list <arg1> <arg2>)` → `(<arg1> <arg2>)`
- **Quoting** does the same thing... but expressions that are not self-evaluating (i.e. variables or procedure calls) will not be evaluated.

The difference between `list` and `'`

```
scm> (define a 1)
```

```
a
```

```
scm> (define b 2)
```

```
b
```

```
scm> (list a b)
```

```
(1 2)
```

```
scm> '(a b)
```

```
(a b)
```

List examples

```
scm> (equal? '(1 2) (list 1 2))
```

```
true
```

```
scm> '(1 . (2 3))
```

```
(1 2 3)
```

```
scm> '(define (square x) (* x x))
```

```
(define (square x) (* x x))
```

append

A useful procedure for **concatenating lists** that never seems to officially get covered. Takes in zero or more **lists** (not list **elements!**), and returns a single well-formed list containing all the elements of the input lists, in order.

```
scm> (append <lst1> <lst2> ...)  
(<lst1 elements> <lst2 elements> ...)
```

If you pass in no arguments, it returns `nil`. It also has robust behavior for random `nils` as arguments:

```
scm> (append nil '(1 (2)) nil '(3) nil nil '(5))  
(1 (2) 3 5)
```


thanks for coming everyone!

:)